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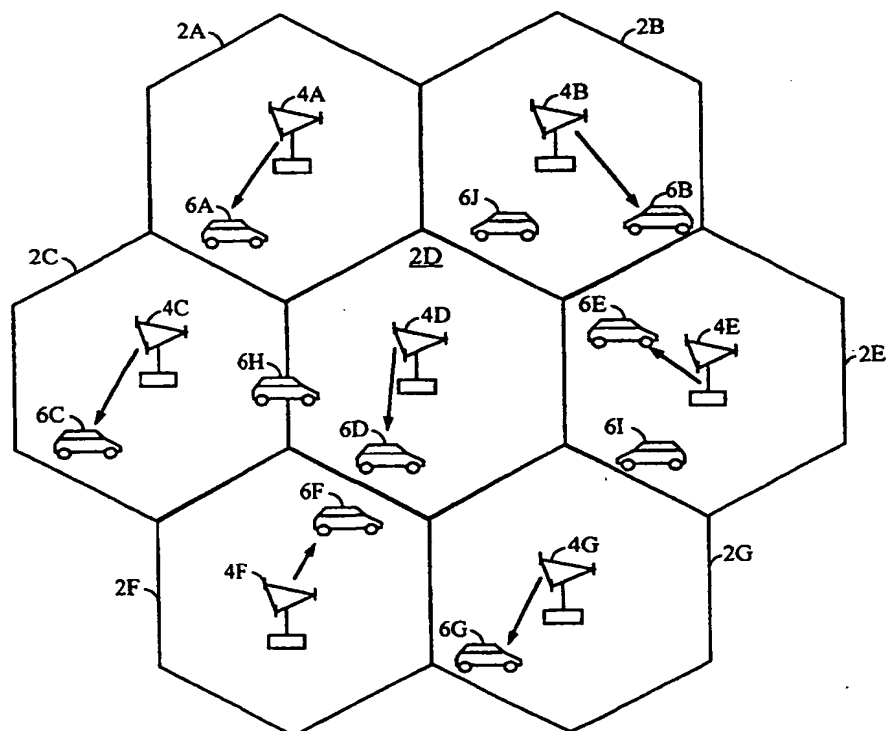
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(54) Title: **SYSTEM FOR ALLOCATING RESOURCES IN A COMMUNICATION SYSTEM**

(57) Abstract

A communication network having a plurality of subscriber units (6) receive a finite resource from a common node is disclosed. Individual subscriber units (6) may seize the finite resource of the common node to the exclusion of all other subscriber units (6) in the network. A scheduler (12) allocates the finite resource to the individual subscriber units (6) based upon a weight associated with the individual subscriber units (6). The scheduler (12) determines the weight for each of the subscriber units (6) based upon an instantaneous rate of consuming the finite resource.



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SYSTEM FOR ALLOCATING RESOURCES IN A COMMUNICATION SYSTEM

BACKGROUND OF THE INVENTION

I. Field of the Invention:

Embodiments disclosed herein relate to communication systems. Particularly, these embodiments are directed to allocating communication resources among the plurality of subscribers to a communication system.

II. Related Art:

Several solutions have been presented to address the problem of allocating limited communication resources provided by a single node in a communication system among a plurality of subscribers. It is an objective of such systems to provide sufficient resources at the nodes to satisfy the requirements of all subscribers while minimizing costs. Accordingly, such systems are typically designed with the objective of efficient allocation of resources among the various subscribers.

Various systems have implemented a frequency division multiple access (FDMA) scheme which allocates resources to each of the subscribers concurrently. A communication node in such systems typically has a limited bandwidth for either transmitting information to or receiving information from each subscriber in the network at any point in time. This scheme typically involves allocating distinct portions of the total bandwidth to the individual subscribers. While such a scheme may be effective for systems in which subscribers require uninterrupted communication with the communication node, better utilization of the total bandwidth may be achieved when such constant, uninterrupted communication is not required.

Other schemes for allocating communication resources of a single communication node among a plurality of subscribers includes time division multiple access (TDMA) schemes. These TDMA schemes are particularly effective in allocating the limited bandwidth resources of a single communication node among a plurality of subscribers which do not require constant, uninterrupted communication with the single communication node. TDMA schemes typically dedicate the entire bandwidth of the single communication node to each of the subscribers at designated time intervals. In a wireless communication system which

employs a code division multiple access (CDMA) scheme, this may be accomplished by assigning to each of the subscriber units all code channels at the designated time intervals on a time multiplexed basis. The communication node implements the unique carrier frequency or channel
5 code associated with the subscriber to enable exclusive communication with the subscriber. TDMA schemes may also be implemented in land line systems using physical contact relay switching or packet switching.

TDMA systems typically allocate equal time intervals to each subscriber in a round robin fashion. This may result in an under utilization
10 of certain time intervals by certain subscribers. Similarly, other subscribers may have communication resource requirements which exceed the allocated time interval, leaving these subscribers under served. The system operator then has the choice of either incurring the cost of increasing the bandwidth of the node to ensure that none of the subscribers are under
15 served, or allowing the under served subscribers to continue to be under served.

Accordingly, there is a need to provide a system and method of allocating communication resources among subscribers to a communication network efficiently and fairly according to a network policy of allocating the
20 communication resources among the subscribers.

SUMMARY OF THE INVENTION

An object of an embodiment of the present invention is to provide a
25 system and method for allocating a finite resource of a communication system among a plurality of subscribers.

Another object of an embodiment of the present invention is to provide a system and method for allocating data transmission resources among a plurality of subscribers which have varying capacities to receive
30 data.

It is another object of an embodiment of the present invention to provide a system and method for optimally allocating data transmission resources among a plurality of subscribers subject to a fairness criteria according to a network policy.

35 It is another object of an embodiment of the present invention to provide a system and method for allocating data transmission resources of a base station among a plurality of remote stations in a wireless communication network.

It is yet another object of an embodiment of the present invention to provide a system and method for enhancing the efficiency of transmitting data to a plurality of subscribers in a variable rate data transmission network by allocating transmission resources to each individual subscriber based upon the rate at which the subscriber can receive transmitted data.

Briefly, an embodiment of the present invention is directed to a resource scheduler in a communication system which includes a common node and a plurality of customer nodes associated with the common node. The common node, at any particular service interval, is capable of providing a finite resource to be seized by one or more engaging customer nodes to the exclusion of any remaining customer nodes. The resource scheduler includes logic for maintaining a weight or score associated with each of the customer nodes, logic for selecting one or more of the remaining customer nodes to seize the finite resource in a subsequent service interval based upon a comparison of the weight associated with each of the selected customer nodes and the respective weights associated with the other remaining customer nodes, and logic for changing the weights associated with the customer nodes to cause an optimal allocation of the finite resource subject to a fairness criteria.

The resource scheduler may maintain the weights associated with each customer node based upon the instantaneous rate at which the customer node can receive data from the common node. The resource scheduler may then favor transmission to the customer nodes having the higher rates of receiving data. By maintaining a weight associated with each of the customer nodes, and selecting individual customer nodes to seize the common node, the scheduler can optimally allocate resources to the customer nodes subject to a fairness criteria.

In the embodiment where the common node provides data transmission resources to the customer nodes, for example, the scheduler may apply weights to the individual customer nodes so as to favor those customer nodes capable of receiving data at higher rates. Such a weighting tends to enhance the overall data throughput of the common node. In another embodiment, the weights are applied in a manner so that the scheduler also complies with the fairness criteria.

While the embodiments disclosed herein are directed to methods and systems for allocating data transmission resources to subscribers through a forward channel in a data service network, the underlying principles have even broader applications to the allocation of resources among elements in a communication system generally. The disclosed embodiments are

therefore intended to be exemplary and not limiting the scope of the claims. For example, principles described herein are applicable to communication networks in which the customer nodes compete for the ability to transmit data to a common node through a limited reverse transmission channel.

5

BRIEF DESCRIPTION OF THE FIGURES

Figure 1 shows a communication network according to an embodiment of the present invention.

10 Figure 2 shows a schematic diagram illustrating details of an embodiment of a base station controller in the communication network illustrated in Figure 1.

Figure 3 shows a flow diagram illustrating the execution of a scheduling algorithm in an embodiment of the channel scheduler shown in
15 Figure 2.

Figure 4 shows a diagram illustrating the timing of the execution of an embodiment of the scheduling algorithm shown in Figure 3.

Figure 5 shows flow diagram illustrating an embodiment of the process for updating the weights for a selected queue in the embodiment
20 identified in Figure 3.

Figures 6a through 6c show a flow diagram illustrating a first embodiment of the process for selecting a queue to receive data transmission in a service interval identified in Figure 3.

Figures 7a through 7d show a flow diagram illustrating a second
25 embodiment of the process for selecting a queue to receive data transmission in a service interval identified in Figure 3.

Figures 8a and 8b show a flow diagram illustrating a third embodiment of the process for selecting a queue to receive data transmission in a service interval identified in Figure 3.

30

DETAILED DESCRIPTION

Embodiments of the present invention are directed to a system and apparatus for allocating resources among a plurality of subscribers to a
35 communication network which are serviced by a single communication node. At individual discrete transmission intervals, or "service intervals," individual subscribers seize a finite resource of the communication node to the exclusion of all other subscribers. The individual subscribers are selected to seize the finite resource based upon a weight or score associated with the

individual subscribers. Changes in a weight associated with an individual subscriber are preferably based upon an instantaneous rate at which the individual subscriber is capable of consuming the finite resource.

Referring to the figures, Figure 1 represents an exemplary variable rate communication system. One such system is described in the U.S. Patent Application Serial No. 08/963,386, entitled Method and Apparatus for High Rate Packet Data Transmission, filed on November 3, 1997, assigned to Qualcomm, Inc. and incorporated herein by reference. The variable rate communication system comprises multiple cells 2a-2g. Each cell 2 is serviced by a corresponding base station 4. Various remote stations 6 are dispersed throughout the communication system. In the exemplary embodiment, each of remote stations 6 communicates with at most one base station 4 on a forward link at any data transmission interval. For example, base station 4a transmits data exclusively to remote station 6a, base station 4b transmits data exclusively to remote station 6b, and base station 4c transmits data exclusively to remote station 6c on the forward link at time slot n. As shown by Figure 1, each base station 4 preferably transmits data to one remote station 6 at any given moment. In other embodiments, the base station 4 may communicate with more than one remote station 6 at a particular data transmission interval to the exclusion of all other remote stations 6 associated with the base station 4. In addition, the data rate is variable and is dependent on the carrier-to-interference ratio (C/I) as measured by the receiving remote station 6 and the required energy-per-bit-to-noise ratio (E_b/N_0). The reverse link from remote stations 6 to base stations 4 is not shown in Figure 1 for simplicity. According to an embodiment, the remote stations 6 are mobile units with wireless transceivers operated by wireless data service subscribers.

A block diagram illustrating the basic subsystems of an exemplary variable rate communication system is shown in Figure 2. Base station controller 10 interfaces with packet network interface 24, public switched telephone network (PSTN) 30, and all base stations 4 in the communication system (only one base station 4 is shown in Figure 2 for simplicity). Base station controller 10 coordinates the communication between remote stations 6 in the communication system and other users connected to packet network interface 24 and PSTN 30. PSTN 30 interfaces with users through a standard telephone network (not shown in Figure 2).

Base station controller 10 contains many selector elements 14, although only one is shown in Figure 2 for simplicity. Each selector element 14 is assigned to control communication between one or more base

stations 4 and one remote station 6. If selector element 14 has not been assigned to remote station 6, call control processor 16 is informed of the need to page remote station 6. Call control processor 16 then directs base station 4 to page remote station 6.

5 Data source 20 contains a quantity of data which is to be transmitted to the remote station 6. Data source 20 provides the data to packet network interface 24. Packet network interface 24 receives the data and routes the data to the selector element 14. Selector element 14 transmits the data to each base station 4 in communication with remote station 6. In the
10 exemplary embodiment, each base station 4 maintains a data queue 40 which stores the data to be transmitted to the remote station 6.

 The data is transmitted in data packets from data queue 40 to channel element 42. In the exemplary embodiment, on the forward link, a "data packet" refers to a quantity of data which is the maximum of 1024 bits and a
15 quantity of data to be transmitted to a destination remote station 6 within a "time slot" (such as ≈ 1.667 msec). For each data packet, channel element 42 inserts the necessary control fields. In the exemplary embodiment, channel element 42 CRC encodes the data packet and control fields and inserts a set
20 of code tail bits. The data packet, control fields, CRC parity bits, and code tail bits comprise a formatted packet. In the exemplary embodiment, channel element 42 then encodes the formatted packet and interleaves (or reorders) the symbols within the encoded packet. In the exemplary embodiment, the interleaved packet is covered with a Walsh code, and spread with the short PNI and PNQ codes. The spread data is provided to RF unit 44 which
25 quadrature modulates, filters, and amplifies the signal. The forward link signal is transmitted over the air through antenna 46 on forward link 50.

 At remote station 6, the forward link signal is received by antenna 60 and routed to a receiver within front end 62. The receiver filters, amplifies, quadrature demodulates, and quantizes the signal. The digitized signal is
30 provided to demodulator (DEMOD) 64 where it is despread with the short PNI and PNQ codes and discovered with the Walsh cover. The demodulated data is provided to decoder 66 which performs the inverse of the signal processing functions done at base station 4, specifically the de-interleaving, decoding, and CRC check functions. The decoded data is
35 provided to data sink 68.

 The hardware, as pointed out above, supports variable rate transmissions of data, messaging, voice, video, and other communications over the forward link. The rate of data transmitted from the data queue 40 varies to accommodate changes in signal strength and the noise

environment at the remote station 6. Each of the remote stations 6 preferably transmits a data rate control (DRC) signal to an associated base station 4 at each time slot. The DRC signal provides information to the base station 4 which includes the identity of the remote station 6 and the rate at which the remote station 6 is to receive data from its associated data queue. Accordingly, circuitry at the remote station 6 measures the signal strength and estimates the noise environment at the remote station 6 to determine the rate at which information which is to be transmitted in the DRC signal.

Embodiments of the present invention are applicable to other hardware architectures which can support variable rate transmissions. The reverse link is not shown nor described for simplicity. However, the present invention can be readily extended to cover variable rate transmissions on the reverse link. For example, instead of determining the rate of receiving data at the base station 4 based upon a DRC signal from remote stations 6, the base station 4 measures the strength of the signal received from the remote stations 6 and estimates the noise environment to determine a rate of receiving data from the remote station 6. The base station 4 then transmits to each associated remote station 6 the rate at which data is to be transmitted in the reverse link from the remote station 6. The base station 4 may then schedule transmissions on the reverse link based upon the different data rates on the reverse link in a manner similar to that described herein for the forward link.

Also, a base station 4 of the embodiment discussed above transmits to a selected one, or selected ones, of the remote stations 6 to the exclusion of the remaining remote stations associated with the base station 4 using a code division multiple access (CDMA) scheme. At any particular time, the base station 4 transmits to the selected one, or selected ones, of the remote station 6 by using a code which is assigned to the receiving base station(s) 4. However, the present invention is also applicable to other systems employing different time division multiple access (TDMA) methods for providing data to select base station(s) 4, to the exclusion of the other base stations 4, for allocating transmission resources optimally.

The channel scheduler 12 connects to all selector elements 14 within the base station controller 10. The channel scheduler 12 schedules the variable rate transmissions on the forward link. The channel scheduler 12 receives the queue size, which is indicative of the amount of data to transmit to remote station 6, and messages from remote stations 6. The channel scheduler 12 preferably schedules data transmissions to achieve the

system goal of maximum data throughput while conforming to fairness a constraint.

As shown in Figure 1, remote stations 6 are dispersed throughout the communication system and can be in communication with zero or one base station 4 on the forward link. In the exemplary embodiment, channel scheduler 12 coordinates the forward link data transmissions over the entire communication system. A scheduling method and apparatus for high speed data transmission are described in detail in U.S. Patent Application Ser. No. 08/798,951, entitled "Method and Apparatus for Forward Link Rate Scheduling," filed February 11, 1997, assigned to the assignee of the present invention and incorporated by reference herein.

According to an embodiment, the channel scheduler 12 is implemented in a computer system which includes a processor, random access memory (RAM) and a program memory for storing instructions to be executed by the processor (not shown). The processor, RAM and program memory may be dedicated to the functions of the channel scheduler 12. In other embodiments, the processor, RAM and program memory may be part of a shared computing resource for performing additional functions at the base station controller 10. In the present embodiment, an individual channel scheduler 12 is distributed to each of the base stations 4. In other embodiments, a single channel scheduler may be centralized for scheduling the transmissions for all base stations 4.

Figure 3 shows an embodiment of a scheduling algorithm which controls the channel scheduler 12 to schedule transmissions from the base station 4 to the remote stations 6. As discussed above, a data queue 40 is associated with each remote station 6. The channel scheduler 12 associates each of the data queues 40 with a "weight" which is evaluated at a step 110 for selecting the particular remote station 6 associated with the base station 4 to receive data in a subsequent service interval. The channel scheduler 12 selects individual remote stations 6 to receive a data transmission in discrete service intervals. At step 102, the channel scheduler initializes the weight for each queue associated with the base station 4.

A channel scheduler 12 cycles through steps 104 through 112 at transmission intervals or service intervals. At step 104, the scheduler 12 determines whether there are any additional queues to be added due to the association of an additional remote station 6 with the base station 4 detected in the previous service interval. The channel scheduler 12 also initializes the weights associated with the new queues at step 104. As discussed above, the base station 4 receives the DRC signal from

each remote station 6 associated therewith at regular intervals such as time slots.

This DRC signal also provides the information which the channel scheduler uses at step 106 to determine the instantaneous rate for consuming information (or receiving transmitted data) for each of the remote stations associated with each queue. According to an embodiment, a DRC signal transmitted from any remote station 6 indicates that the remote station 6 is capable of receiving data at any one of eleven effective data rates shown in Table 1. Such a variable rate transmission system is described in detail in U.S. Patent Application Serial No. 08/965,705 referenced above.

Table 1

Effective Rate (R_i)	Data Transmitted in Service Interval (Data_Size (L_i)) (bits)	Length/Transmission Time of Service Interval (L_i) (time slots \approx 1.667 msec)
38.4 kbps	1024	16
76.8 kbps	1024	8
102.4 kbps	1024	6
153.6 kbps	1024	4
204.8 kbps	1024	3
307.2 kbps	1024	2
614.4 kbps	1024	1
921.6 kbps	1536	1
1228.8 kbps	2048	1
1843.2 kbps	3072	1
2457.6 kbps	4096	1

The channel scheduler at step 108 determines the length of a service interval during which data is to be transmitted to any particular remote station based upon the remote station's associated instantaneous rate for

receiving data (as indicated in the most recently received DRC signal). According to an embodiment, the instantaneous rate of receiving data R_i determines the service interval length L_i associated with a particular data queue at step 106. Table 1 summarizes the L_i values for each of the eleven possible rates for receiving data at a remote station 6.

The channel scheduler at step 110 selects the particular data queue for transmission. The associated quantity of data to be transmitted is then retrieved from a data queue 40 and then provided to the channel element 42 for transmission to the remote station 6 associated with the data queue 40. As discussed below, the channel scheduler 12 at step 110 selects the queue for providing the data which is transmitted in a following service interval using information including the weight associated with each of the queues. The weight associated with the transmitted queue is then updated at step 112.

Figure 4 shows a diagram illustrating the timing of the scheduler 12 and data transmission in service intervals. Figure 4 shows three discrete service intervals during transmission at time intervals δ_{-1} , δ_0 and δ_1 . As steps 104 through 112 of the scheduling algorithm of Figure 4 are executed during service intervals 202, the scheduling algorithm executing during the interval δ_0 preferably determines which queue is to be transmitted at the interval δ_1 . Also, as discussed below, the execution of steps 104 through 112 relies on information in the DRC signals received from the remote stations 6. This information is preferably extracted from the most recently received DRC signals. Accordingly, the steps 104 through 110 are preferably executed and completed during the last time slot of the service intervals. This ensures that the decisions for allocating the subsequent service interval are based upon the most recent DRC signals (i.e., those DRC signals which are in the time slot immediately preceding the execution of the steps 104 through 110).

Steps 104 and 110 are preferably completed within a time slot while providing sufficient time the channel scheduler 12 to schedule the transmissions for the subsequent service interval. Thus the processor and RAM employed in the channel scheduler 12 are preferably capable of performing the steps 104 through 112 within the time constraints illustrated in Figure 4. That is, the processor and RAM are preferably sufficient to execute steps 104 through 110, starting at the beginning of a time slot and completing steps 104 through 110, within sufficient time before the end of the time slot for the channel scheduler 12 to schedule transmissions in a subsequent service interval.

Figure 5 shows an embodiment of the process for updating the weights at step 112 (Figure 3). Step 302 computes a rate threshold "C" which is an average of all of the instantaneous rates associated with queues having data. The instantaneous rates associated with queues which do not include data are preferably eliminated for this calculation. Step 304 compares the instantaneous rate associated with the Selected _Queue selected at step 110. If an instantaneous rate associated with a Selected _Queue exceeds the threshold C, step 306 increments the weight associated with this Selected _Queue by a lower value which is preferably a number representing the quantity of data to be transmitted during the subsequent service interval from the Selected _Queue in units such as bits, bytes or megabytes. If the instantaneous rate associated with the Selected _Queue does not exceed the threshold calculated at step 302, step 308 increments the weight of the Selected _Queue by a higher value which is preferably a multiple "G" of the quantity of data which is to be transmitted during the subsequent service interval from the Selected _Queue such as bits, bytes or megabyte quantities.

The selection of G is preferably based upon a fairness criteria which favors the allocation of service intervals to remote stations 6 having the capacity to receive data at higher rates. The system designer selects the size of G based upon the extent to which remote stations 6 receiving data at the higher rates are to be favored over the slower receiving remote stations 6. The larger the value of G, the more efficiently the forward link of the base station 4 is utilized. This efficiency, however, comes at the cost of depriving the subscribers of the slower receiving remote station 6 of the transmission resources of the forward link. The system designer therefore preferably selects the value of G in a manner which balances the two competing objectives of 1) enhancing the overall efficiency of the forward link and 2) preventing accute deprivation of the slower receiving remote stations 6.

Steps 304, 306 and 308 illustrate that selected queues having a faster associated instantaneous data rate (i.e., exceeding the threshold C) will tend to have the associated weight incremented by only a small amount, while selected queues having a lower data rate (i.e., not exceeding the threshold C) will have its associated weight incremented by a significantly greater amount. As discussed below in connection with the algorithm performed at step 110 of Figure 3, this implementation tends to favor servicing remote stations which receive data at relatively faster rates over those remote stations receiving data at lower data rates.

This tendency enhances the throughput efficiency of the base station 4 in transmitting data in the forward link. However, as the weights associated

with the often selected queues associated with the remote stations having the higher rates of receiving data (i.e., exceeding the threshold C) continue to be incremented, these weights eventually approach the weights of the queues associated with the less often selected queues associated with the remote stations having the slower rates of receiving data (i.e., not exceeding the threshold). The selection process at step 110 will then begin to favor the slower receiving remote stations as the weights of the faster receiving remote stations begin to exceed the weights of the slower receiving remote stations. This imposes a fairness restraint on the selection process at step 110 by preventing the faster receiving remote stations from dominating the forward link transmission resources of the base station to the exclusion of the slower receiving remote stations.

It is an objective of the present embodiment to ensure that queues having no data to transmit are not given an unfair preference for transmission over those queues having data. At steps 102 and 104, all new queues are initialized with a weight of zero. Without being selected, such queues will continue to maintain the weight of zero provided that the queue is not selected. Therefore, step 310 in Figure 5 decrements the weight of all queues, to a value no less than zero, by the minimum weight of any queue with data (determined at step 309). This is illustrated in detail below in an example shown in Table 2.

Table 2

Service Interval	Weights at the End of the Service Interval			Remote Station Selected in Service Interval	Remote Station Serviced in Service Interval	Amount by Which Weights are Decrementd
	Remote Station 1	Remote Station 2	Remote Station 3			
0	0	0	0	N/A	N/A	N/A
1	1	0	0	1	N/A	0
2	1	1	0	2	1	0
3	0	0	7	3	2	1
4	1	0	7	1	3	0
5	0	0	6	2	1	1
6	1	0	6	1	2	0
7	0	0	5	2	1	1

This example has three remote stations each associated with a queue of data to be transmitted from a base station. The example assumes that remote station 1 has the highest data rate, remote station 2 has the next highest data rate and remote station 3 has the lowest data rate. For simplicity, it is assumed that these data rates do not change over the service intervals 1 through 7. It is also assumed that the data rates at remote station 1 and remote station 2 each exceed the threshold C at step 304, and that the data rate associated with remote station 3 does not exceed this threshold. It is further assumed that step 306 will increment the weight of the Selected_Queue by one if the Selected_Queue is associated with the remote station 1 or remote station 2, and that step 308 will increment the weight of the Selected_Queue by eight if the Selected_Queue is associated with the remote station 3.

At service interval 1, the channel scheduler 12 selects the remote station 1 to receive data in the subsequent service interval, since, while it has the lowest weight along with remote stations 2 and 3, remote station 1 has a higher rate of receiving data. Data is then transmitted to remote station 1 during service interval 2 and the weight associated with the remote station 1 is incremented by one at the end of service interval 1. The channel scheduler 12 then selects remote station 2 to receive data in service interval 3 (since remote station 2 has the lowest weight and a faster rate of receiving data than does remote station 3). As shown in Table 2, the weight of remote station 2 is incremented by 1 by the end of the service interval 2.

At the beginning of service interval 3, remote station 3 has the lowest weight. The channel scheduler 12 selects remote station 3 to receive data at the service interval 4. The state at the end of interval 3 reflects that weight of the remote station 3 was incremented from zero to eight to reflect the selection of the remote station 3. The weights at the remote stations 1, 2 and 3 are then decremented by one which is consistent with step 310 (Figure 5) as indicated in Table 2. At service interval 4, the channel scheduler 12 selects remote station 1 to receive data in service interval 4 since the queue associated with remote station 1 has the lowest weight and the highest rate for receiving data.

The channel scheduler 12 at service interval 5 selects remote station 2 to receive data during service interval 6. The weight associated with the remote station 2 is first incremented at step 306 and the weights of all of the remote stations are decremented by one as reflected in the weights at the end of the service interval 5 as shown in Table 2. Remote station 1, having the

lowest weight, is then selected again in service interval 6 for receiving data in service interval 7.

As shown in the embodiment of Figure 1, the remote stations 6 are mobile and capable of changing associations among the different base stations 4. For example, a remote station 6f is initially receiving data transmissions from the base station 4f. The remote station 6f may then move out of the cell of the base station 4f and into the cell of the base station 4g. The remote station 6f can then start transmitting its DRC signal to alert the base station 4g instead of the base station 4f. By not receiving a DRC signal from the remote station 6f, logic at the base station 4f deduces that the remote station 6f has disengaged and is no longer to receive data transmissions. The data queue associated with the remote station 6f may then be transmitted to the base station 4g via a land line or RF communication link.

According to an embodiment of the present invention, the channel scheduler 12 at a base station 4 assigns a weight to a queue of a remote station 6 which has disengaged and re-engaged the base station 4. Rather than simply assigning a weight of zero to the re-engaging remote station 6, the base station 4 preferably assigns a weight which does not give the re-engaging remote station an unfair advantage for receiving data transmissions from the base station 4. In one embodiment, the channel scheduler 12 randomly assigns a weight to the queue of the re-engaging remote station 6 according to a uniform distribution between zero and the highest weight of any queue currently serviced by the channel scheduler 12. In another embodiment, the base station 4 receives the weight of the re-engaging remote station 6 from the last base station associated with the remote station 6 via a land line transmission.

In an alternative embodiment, the channel scheduler 12 gives a re-engaging remote station 6 "partial credit" for having a past association with the base station 4. The channel scheduler 12 determines the number of time slots that the previous service interval spans "n," and maintains a history of the number of time slots "m_i" during the previous service interval that the base station 4 received a DRC from the remote station i. The weight of the queue associated with the remote station i is then decremented at step 310 as follows:

$$W_i = W_i - m_i/n \times W_{\min}$$

where:

W_i = the weight of queue i

W_{\min} = the minimum weight of any queue with data to transmit to a remote station

m_i = the number of time slots during the previous service interval that the base station received a DRC from the remote station i

n = the number of time slots that the previous service interval spans

Figures 6a through 6c show a flow diagram illustrating the logic performed at step 110 (Figure 3) according to an embodiment. Step 402 initializes the identity of the Selected _Queue as being the first data queue having data for transmission to an associated remote station 6. At steps 404 through 422, the channel scheduler 12 determines whether this initial queue or a different data queue having data should be selected for transmission to its associated remote station 6. The Next _Queue is then retrieved at step 406 and step 408 determines whether this Next _Queue has data. If the Next _Queue does not have data, execution returns to step 406 to select a subsequent data queue. Otherwise, if this Next _Queue has data, the identity of the Current _Queue is assigned the Next _Queue. If the weight of the Current _Queue exceeds the weight of the Selected _Queue, step 412 returns execution to step 406 to retrieve a subsequent Next _Queue. Otherwise, step 414 determines whether the weight of the Current _Queue is less than the weight of the Selected _Queue. If the weight of the Current _Queue is less than the weight of the Selected _Queue, step 414 moves execution to step 420 to assign the identity of the Current _Queue to the Selected _Queue.

Otherwise, the logic at steps 412 and 414 dictate that if execution reaches step 416, the weights of the Current _Queue and the Selected _Queue are equal. Step 424 assigns the Current _Queue as the Selected _Queue whether the following conditions are met:

- 1) the instantaneous rate of receiving data associated with the Current _Queue exceeds the instantaneous rate of receiving data associated with the Selected _Queue (step 416); and
- 2) if the service interval assigned to the Current _Queue would exhaust all of the data stored in the Current _Queue, leaving a fractional remainder of data in the service interval assigned to the Current _Queue, such a fractional remainder would not exceed any

such fractional remainder of data in the Selected_Queue in the service interval assigned to the Selected_Queue (steps 418 through 422).

Otherwise, execution returns to step 406 to select the Next _Queue.

5 Figures 7a through 7d show a flow diagram illustrating a second embodiment of the logic performed at the step 110 for selecting a queue for transmission to an associated remote station 6. In this embodiment, it is assumed that each base station 4 periodically transmits a control signal to all associated remote stations 6 having a fixed duration (such as eight to sixteen
10 time slots). According to an embodiment, the base station 4 transmits this control signal once every 400 msec. During this control transmission, no data from any data queue 40 (Figure 2) may be transmitted to an associated remote station 6. An objective of the embodiment shown at Figures 7a and 7b is to select only those data queues which may completely transmit for a
15 service interval having a length determined at step 108 before the beginning of the next control signal transmission.

 Steps 499 through 507 filter all of the queues to determine which queues are candidates for completion before the beginning of the next control signal transmission. Step 499 determines the time "T" until the next
20 control signal transmission by, for example, subtracting the scheduled time of the beginning of the next control signal transmission by the beginning of the next scheduled service interval. Step 501 determines whether the length of service interval associated with each queue determined at step 108 can be transmitted within the time T based upon the instantaneous rate of
25 transmission for the remote unit 6 associated with the queue determined at step 106. According to an embodiment, step 501 compares the service interval length with T. Step 502 then determines whether the Next _Queue includes any data. If the Next _Queue satisfies the conditions at steps 501 and 502, the identity of the Next _Queue is assigned to the Selected _Queue.

30 Steps 504 through 508 examine the remaining data queues to determine the data queues having associated service interval (determined at step 108) which may be completely transmitted prior to the beginning of the next control signal transmission. Upon meeting the criteria set forth at steps 507 and 508, the Current _Queue is assigned as the Next _Queue. Steps 512
35 through 526 then perform a selection process according to queue weights in a manner similar to that discussed above in connection with steps 412 through 426 in Figures 6a through 6c. However, in the embodiment of Figures 7a through 7d, only those data queues having an assigned packet length which may be completed prior to the beginning of the next control

signal transmission may be candidates for selection based upon the associated queue weight.

Figures 8a and 8b show a flow diagram illustrating a third embodiment of the logic executed at step 110 at Figure 3 for selecting a queue for transmission. In this embodiment, subscribers of select remote units 6 are guaranteed a minimum average rate of data transmission. For each such premium remote unit, the channel scheduler 12 maintains a timer which alerts the channel scheduler 12 to schedule a transmission to its premium queue, regardless of the weights associated with the remaining queues. The time interval for the particular timer is determined based upon the average data rates guaranteed to the customer, the service interval assigned to that data queue at step 108 (see center column of Table 1), and any instantaneous data rate for receiving data determined at step 106. Thus, the time interval associated with the premium queue timer is dynamic with respect to these values. According to an embodiment, the timer interval is determined whenever the timer is reset as follows:

$$T_j = \frac{\text{Data_Size}(L_j)}{r_j}$$

20

where:

25

T_j = timer interval for premium queue j
 $\text{Data_Size}(L_j)$ = quantity of data to be transmitted in service interval assigned to the premium queue j

30

r_j = average data transmission rate guaranteed to the premium subscriber associated with the premium queue j

35

The timer is reset at either of two events. The first event initiating a reset of the timer is an expiration of the timer interval. The second event for initiating a reset of the timer is a selection of the associated premium data queue based upon its associated weight in a manner discussed above with reference to Figures 6a through 6c.

Steps 606 through 610 determine whether the Next _Queue is a premium queue entitled to a minimum average rate of receiving data and, if so, whether the timer associated with that premium queue has expired. If the timer has expired, step 612 assigns the identity of the Next _Queue to the

Selected _Queue and execution at step 110 completes. The weight of the selected queue is then updated at step 112 as discussed above. If there are no premium queues with an expired timer, step 614 initiates the selection of the queue for transmission in the subsequent service interval at step 616
5 based upon the weights of the queues in a manner discussed above with references to Figures 6a through 6c. If the queue selected at step 616 is a premium queue having an associated timer, step 618 initiates a reset of the timer associated with the selected queue at step 620.

As outlined above, the timer associated with any particular premium
10 data queue is reset following its selection based upon the associated weight at step 620. The associated timer is also reset when it expires before selection of the data queue. The timer thus alerts the channel scheduler 12 to override the logic directed to selecting data queues based upon weights to ensure that this subscriber is associated with the premium data queues
15 receive a guaranteed minimum average rate of receiving data.

While there has been illustrated and described what are presently considered to be the preferred embodiments of the present invention, it will be understood by those skilled in the art that various other modifications may be made, and equivalents may be substituted, without departing from
20 the true scope of the invention. Additionally, many modifications may be made to adapt a particular situation to the teachings of the present invention without departing from the central inventive concept described herein. Therefore, it is intended that the present invention not be limited to the particular embodiments disclosed, but that the invention include all
25 embodiments falling within the scope of the appended claims.

WHAT IS CLAIMED IS:

CLAIMS

1. A resource scheduler in a communication system, the communication system including a common node and a plurality of customer nodes associated with the common node, the common node having at any particular service interval a finite resource seized by one or more engaging customer nodes to the exclusion of any remaining customer nodes, each of the customer nodes having an instantaneous rate of consuming the finite resource, the resource scheduler comprising:

logic for maintaining a weight associated with each of the customer nodes;

logic for selecting one or more of the remaining customer nodes to seize the finite resource in a subsequent service interval based upon a comparison of the weight associated with the selected remaining customer nodes and the respective weights associated with the other remaining customer nodes, the subsequent service interval following a present service interval in which the finite resources are seized by one or more engaging customer nodes; and

logic for changing the weight of the selected customer node based upon the instantaneous rate at which the selected customer node consumes the finite resource.

2. The resource scheduler of claim 1, the resource scheduler further including circuitry for causing the selected one or more customer nodes to engage the common node and seize the finite resource following a termination of the present service interval.

3. The resource scheduler of claim 1, wherein the logic for changing the weight of the selected customer node increments the weight associated with the selected customer node by a value associated with the instantaneous rate at which the selected customer node consumes the finite resource, and wherein the logic for selecting the one or more of the remaining customer nodes selects the one or more remaining customer nodes having one of the lowest weights associated therewith.

4. The resource scheduler of claim 1, wherein the instantaneous rate at which the customer nodes consume the finite resource is dynamic, and wherein the logic for changing the weight of the selected customer node

4 changes the weight of the selected customer node based upon the
instantaneous rate at which the selected customer node consumes the finite
6 resource determined at a particular point in time.

5. The resource scheduler of claim 1, wherein the finite resource
2 includes an instantaneous capacity to transmit information to the selected
ones of the customer nodes, the instantaneous rate of consuming the finite
4 resource is a rate at which the customer service node is capable of receiving
information transmitted from the central node during a service interval,
6 and wherein the common node maintains a memory for storing a data
structure associated with each of the customer nodes having data
8 representative of a quantity of information to be transmitted to the
customer node.

6. The resource scheduler of claim 5, wherein during the
2 subsequent service interval, the common node transmits a quantity of
information to be transmitted to the selected customer which is based upon
4 the rate at which the selected customer node is capable of receiving
information.

7. The resource scheduler of claim 6, wherein the subsequent
2 service interval has an associated subsequent service interval duration for
transmitting data to a selected customer node based upon the instantaneous
4 rate at which the selected customer node is capable of receiving information,
wherein the common node utilizes the finite resource to transmit control
6 information for a control channel duration to at least one of the customer
nodes beginning at fixed intervals, and wherein the logic for selecting the
8 one or more customer nodes selects the one or more customer nodes from
those customer nodes which would result in a subsequent service interval
10 terminating prior to the beginning of a following control channel duration.

8. The resource scheduler of claim 5, wherein the logic for
2 maintaining the weight associated with each of the customer nodes
modifies the weight associated with at least one of the customer nodes when
4 the quantity of information to be transmitted to at least one of the customer
nodes falls below a threshold quantity of information for a specified
6 duration such that the logic for selecting selects from the remaining
customer nodes associated with a quantity of information to be transmitted
8 exceeds the threshold quantity.

9. The resource scheduler of claim 1, wherein the
2 communication system includes a plurality of common nodes, each of the
customer nodes are associated with exactly one of the common nodes at any
4 particular point in time, and at least one of the customer nodes can change
its association between a first common node and a second common node,
6 and wherein the resource scheduler independently maintains the weights
associated with each of the customer nodes associated with at least the first
8 common node, the resource scheduler further including logic for modifying
the weight associated with the at least one customer node based upon a
10 duration of time that the at least one customer node is associated with the
first common node over a specified historical past.

10. The resource scheduler of claim 1, the resource scheduler
2 further including:

logic for determining a duration of an override time interval,
4 the override time interval having a beginning and an end, associated
with at least one customer node based upon a minimum average rate
6 of consuming the finite resource associated with the at least one
customer node and the instantaneous rate of consuming the finite
8 resource associated with the at least one customer node; and

logic for initializing the override time interval whenever the
10 at least one customer node seizes the finite resource and whenever
the override time interval ends,

12 wherein the logic for selecting schedules the at least one
customer node to seize the finite resource in the subsequent service
14 interval independent of the weights associated with the customer
nodes when each override time interval ends.

11. A method of allocating a finite resource in a communication
2 system, the communication system including a common node and a
plurality of customer nodes associated with the common node, the common
4 node having at any particular service interval a finite resource seized by one
or more engaging customer nodes to the exclusion of any remaining
6 customer nodes, each of the customer nodes having an instantaneous rate
of consuming the finite resource, the method comprising:

8 maintaining a weight associated with each of the customer
nodes;

10 selecting one or more of the remaining customer nodes to seize
12 the finite resource in a subsequent service interval based upon a
14 comparison of the weight associated with the selected remaining
16 customer nodes and the respective weights associated with the other
18 remaining customer nodes, the subsequent service interval following
a present service interval in which the finite resources are seized by
one or more engaging customer nodes; and

 changing the weight of the selected customer node based upon
the instantaneous rate at which the selected customer node consumes
the finite resource.

12. The method of claim 11, the method further including causing
the selected one or more customer nodes to engage the common node and
seize the finite resource following a termination of the present service
interval.

13. The method of claim 11, wherein the step of changing the
weight of the selected customer node further includes incrementing the
weight associated with the selected customer node by a value associated with
the instantaneous rate at which the selected customer node consumes the
finite resource, and the step of selecting the one or more of the remaining
customer nodes includes selecting the one or more remaining customer
nodes having one of the lowest weights associated therewith.

14. The method of claim 11, wherein the instantaneous rate at
which the customer nodes consume the finite resource is dynamic, and
wherein the step for changing the weight of the selected customer node
includes changing the weight of the selected customer node based upon the
instantaneous rate at which the selected customer node consumes the finite
resource determined at a particular point in time.

15. The method of claim 11, wherein the finite resource includes
an instantaneous capacity to transmit information to the selected ones of the
customer nodes, the instantaneous rate of consuming the finite resource is a
rate at which the customer service node is capable of receiving information
transmitted from the central node during a service interval, the method
further including maintaining a data structure in a memory at the common
node associated with each customer node, the data structure having data

8 representative of a quantity of information to be transmitted to the
customer node.

2 16. The method of claim 15, the method further including
scheduling a transmission of a quantity of data from the common node to
the selected customer node during the subsequent service interval.

2 17. The method of claim 16, wherein the subsequent service
interval has an associated subsequent service interval duration for
transmitting data to a selected customer node based upon the instantaneous
4 rate at which the selected customer node is capable of receiving information,
wherein the common node utilizes the finite resource to transmit control
6 information for a control channel duration to at least one of the customer
nodes beginning at fixed intervals, and wherein the step for selecting the
8 one or more customer nodes further includes selecting the one or more
customer nodes from those customer nodes which would result in a
10 subsequent service interval terminating prior to the beginning of a
following control channel duration.

2 18. The method of claim 15, wherein the step for maintaining the
weight associated with each of the customer nodes further includes
modifying the weight associated with at least one of the customer nodes
4 when the quantity of information to be transmitted to at least one of the
customer nodes falls below a threshold quantity of information for a
6 specified duration such that the logic for selecting selects from the
remaining customer nodes associated with a quantity of information to be
8 transmitted exceeds the threshold quantity.

2 19. The method of claim 11, wherein the communication system
includes a plurality of common nodes, each of the customer nodes are
associated with exactly one of the common nodes at any particular point in
4 time, and at least one of the customer nodes can change its association
between a first common node and a second common node, and wherein the
6 step for maintaining the weight further includes independently
maintaining the weight associated with each of the customer nodes
8 associated with at least the first common node, the wherein the method
further includes modifying the weight associated with the at least one
10 customer node based upon a duration of time that the at least one customer
node is associated with the first common node over a specified historical

12 past.

20. The method of claim 11, the method further including:

2 determining a duration of an override time interval, the
override time interval having a beginning and an end, associated
4 with at least one customer node based upon a minimum average rate
of consuming the finite resource associated with the at least one
6 customer node and the instantaneous rate of consuming the finite
resource associated with the at least one customer node; and

8 initializing the override time interval whenever the at least
one customer node seizes the finite resource and whenever the
10 override time interval ends,

wherein the step for selecting further includes scheduling the
12 at least one customer node to seize the finite resource in the
subsequent service interval independent of the weights associated
14 with the customer nodes when each override time interval ends.

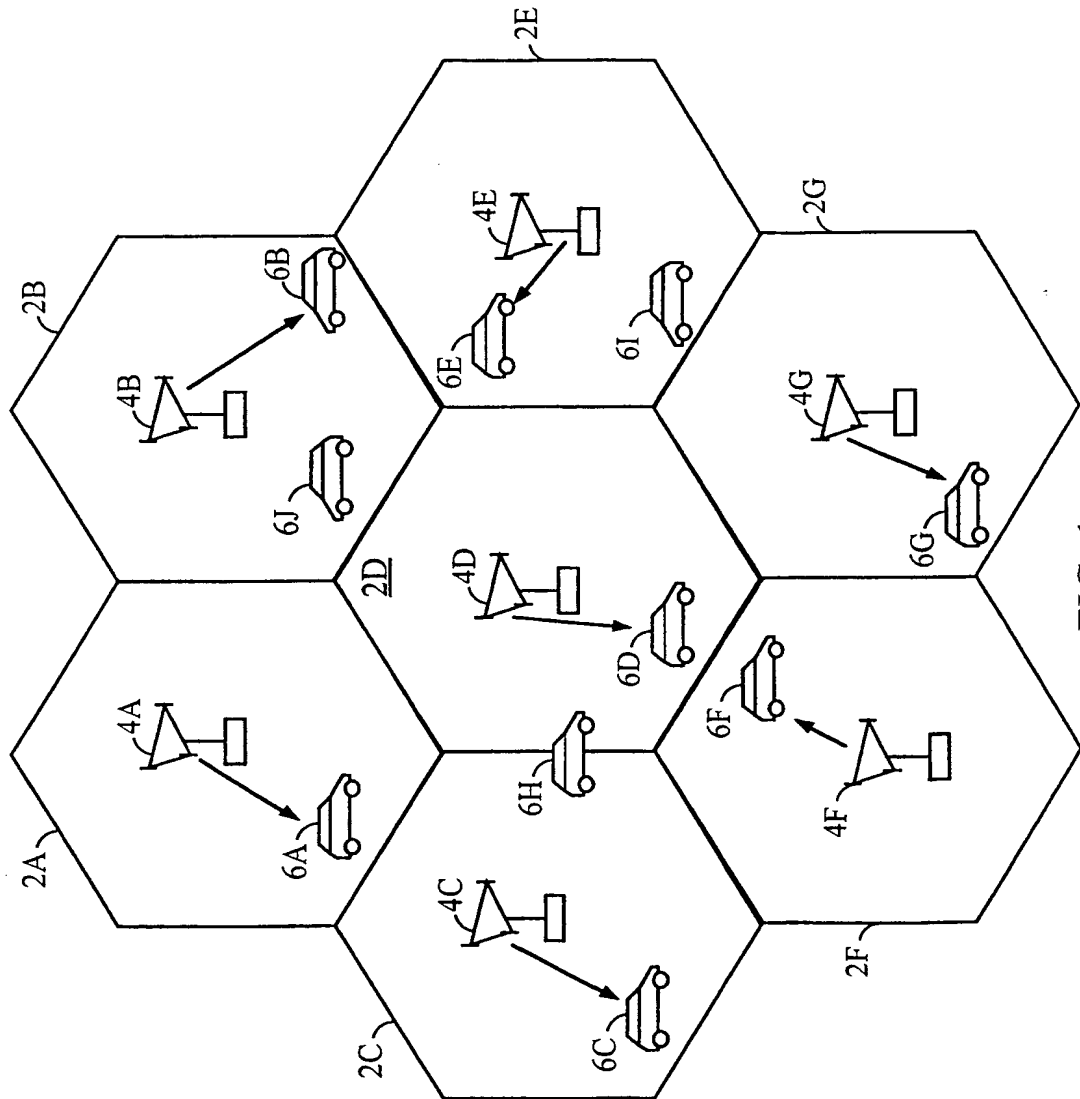


FIG. 1

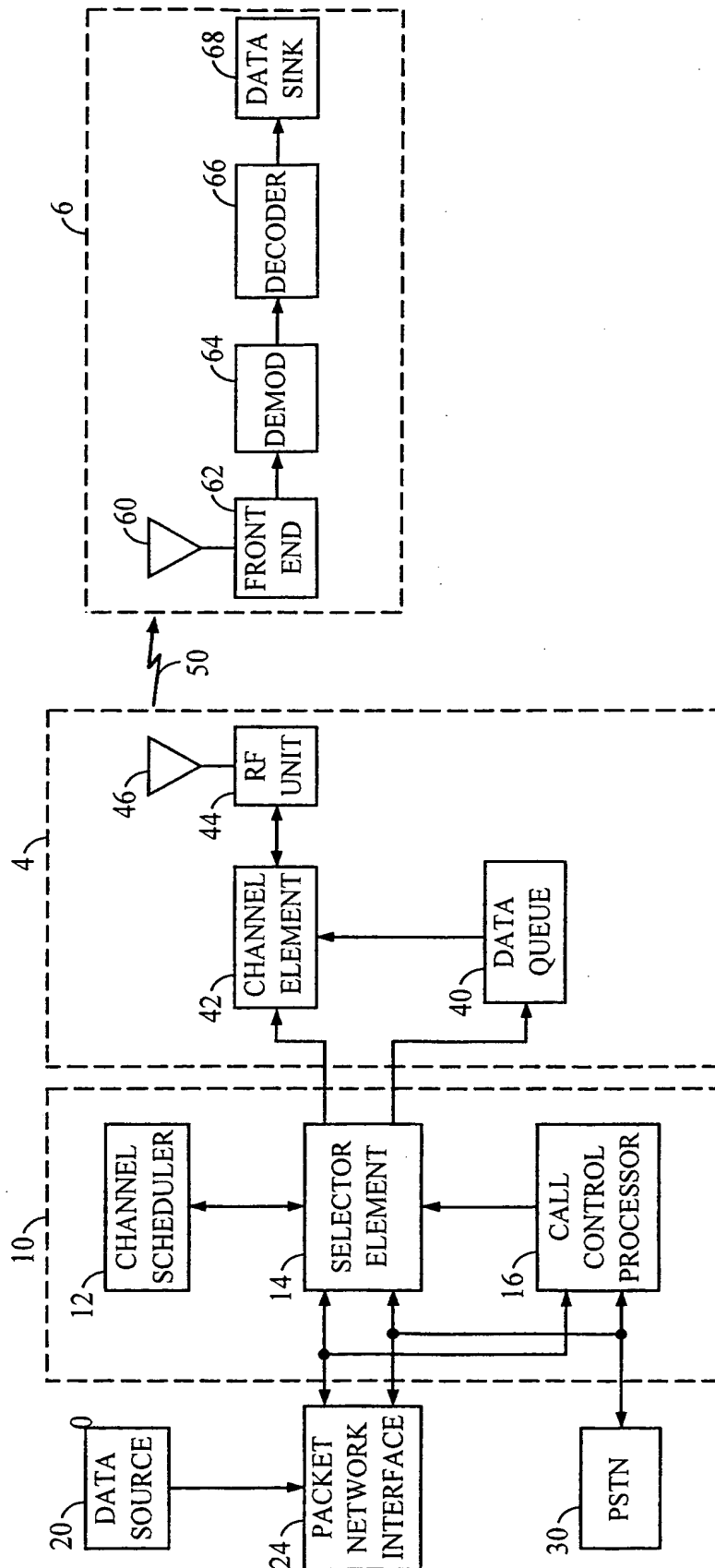


FIG. 2

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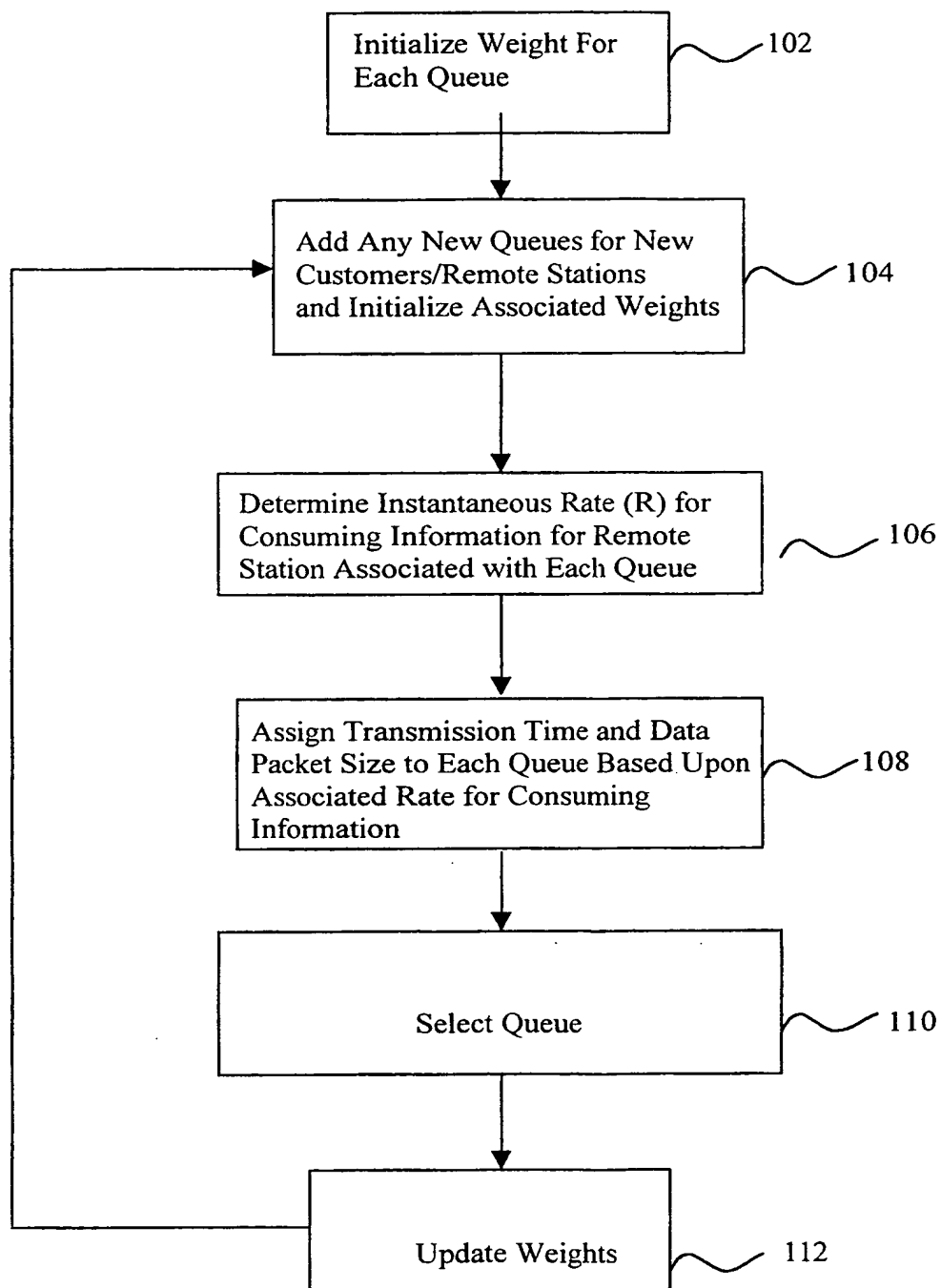


FIGURE 3

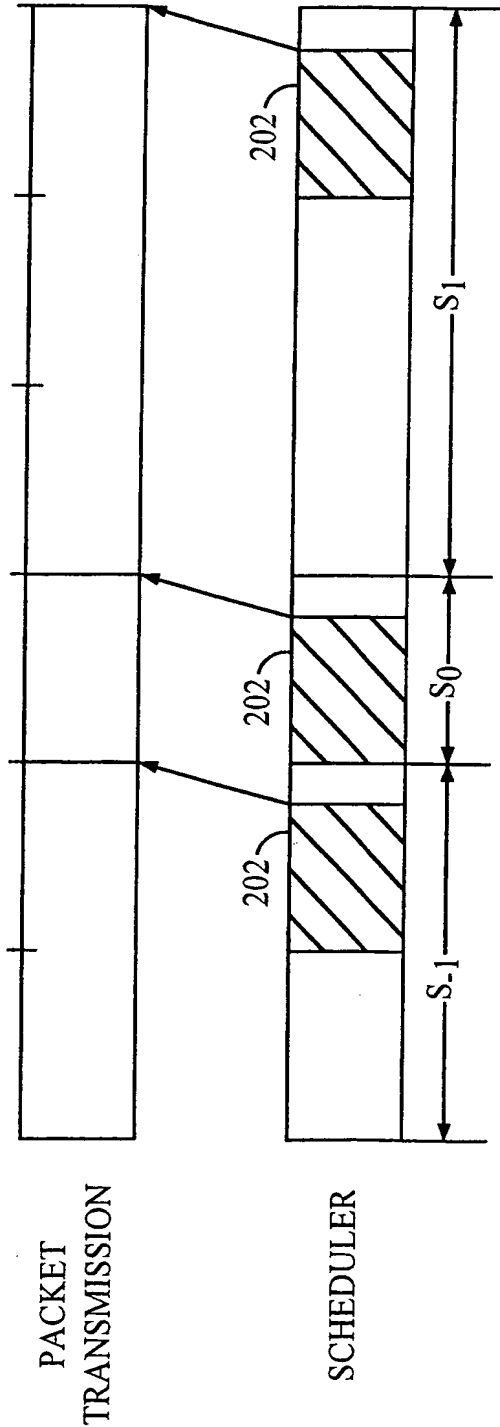


FIG. 4

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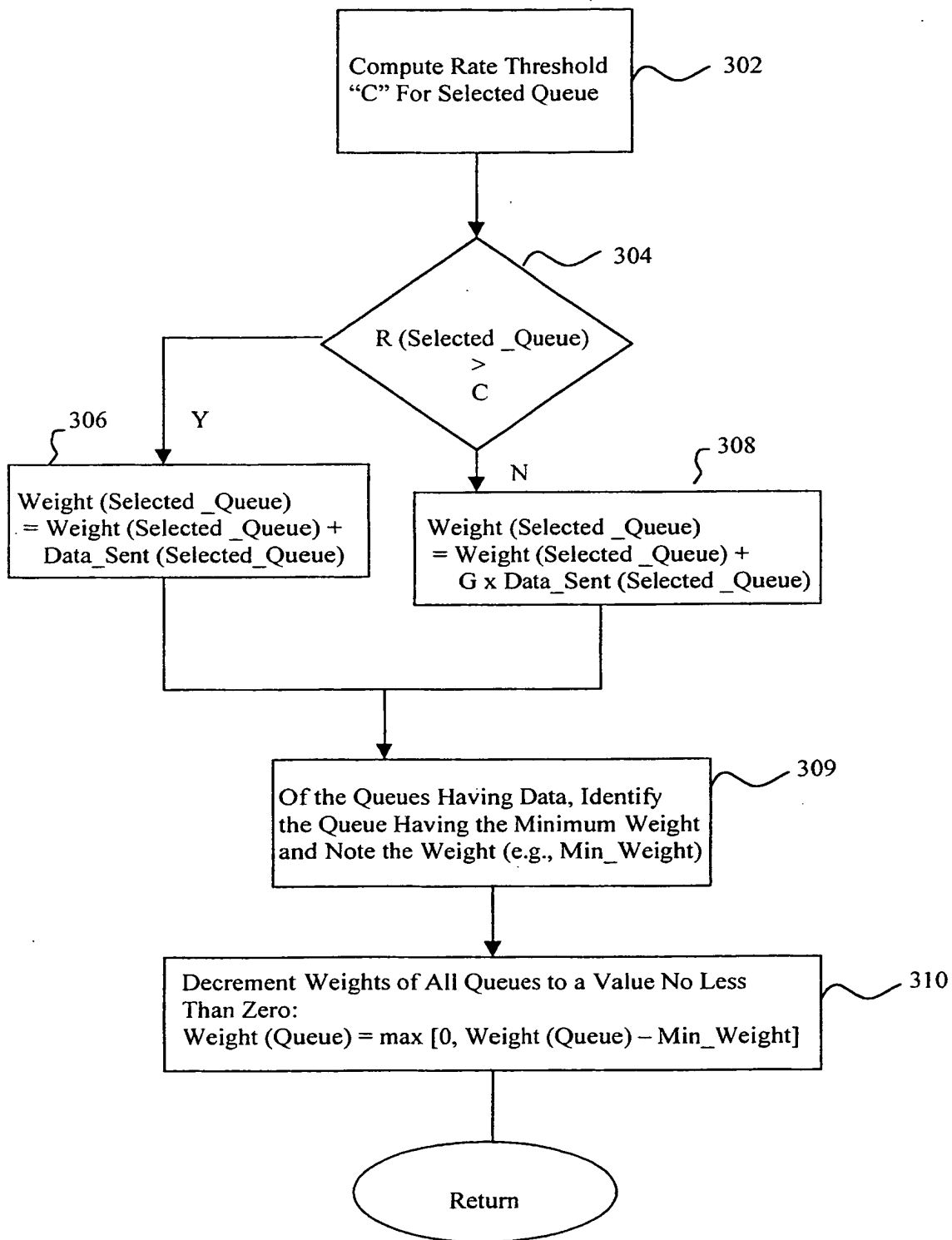


FIGURE 5

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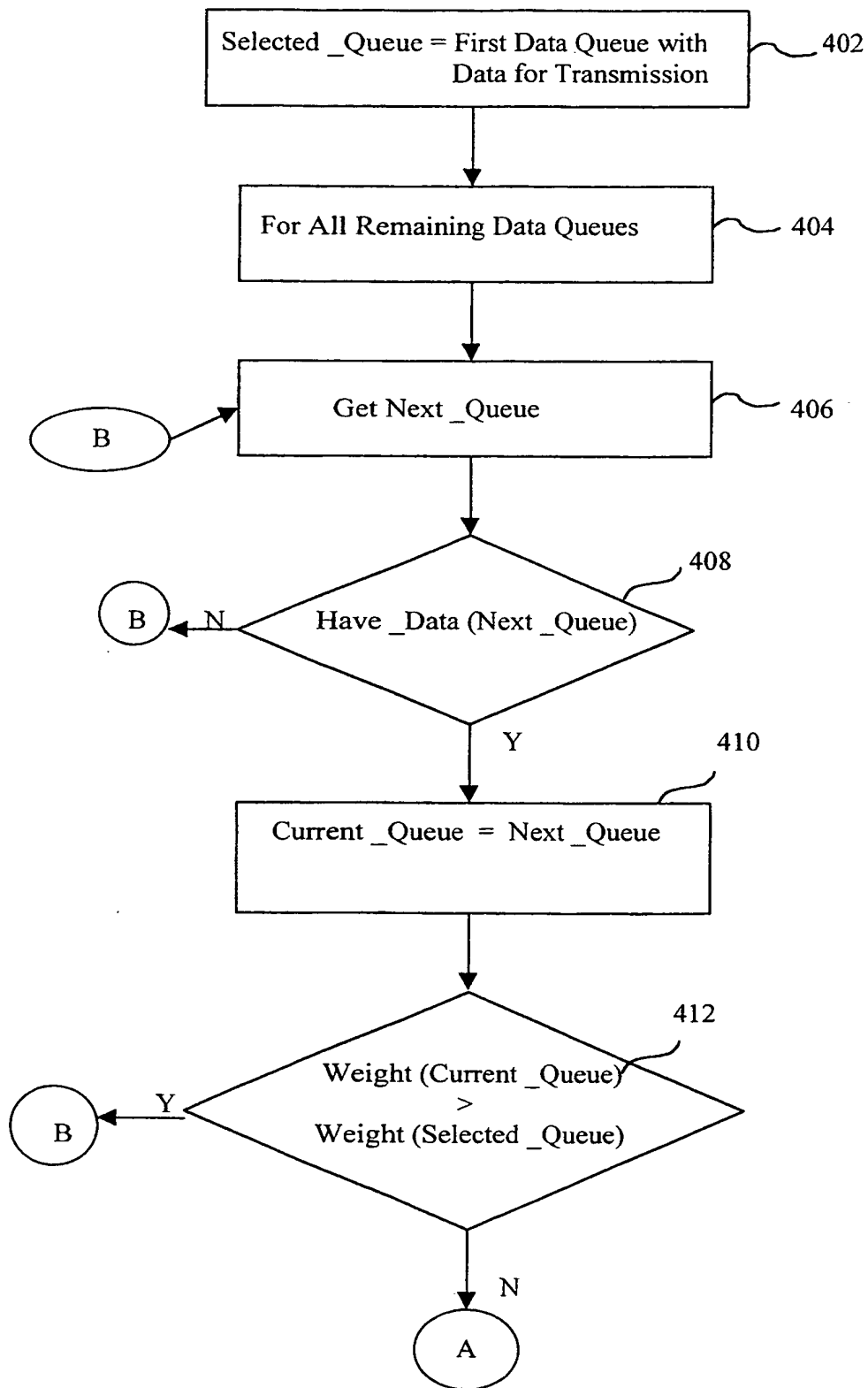


FIGURE 6a

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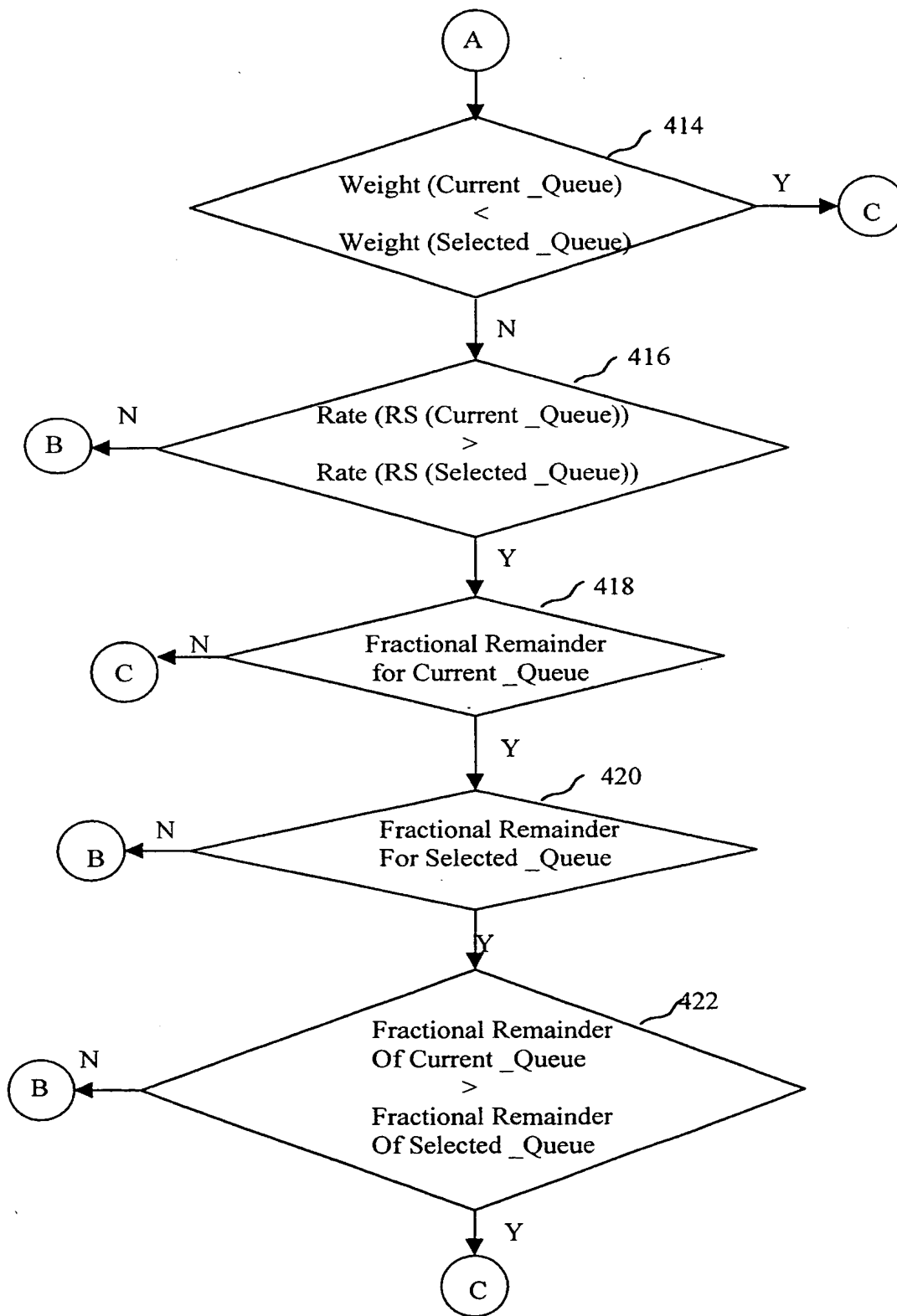


FIGURE 6b

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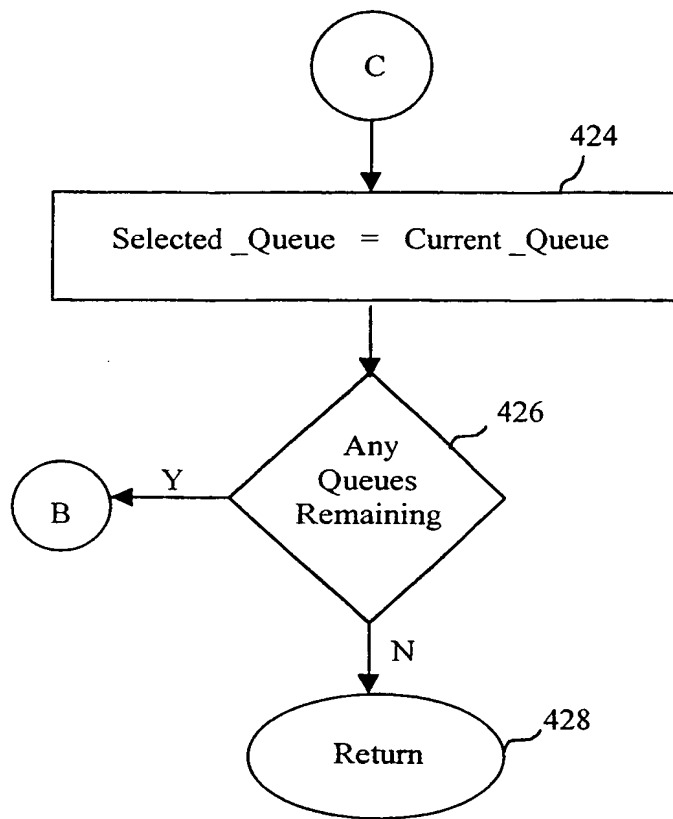


FIGURE 6c

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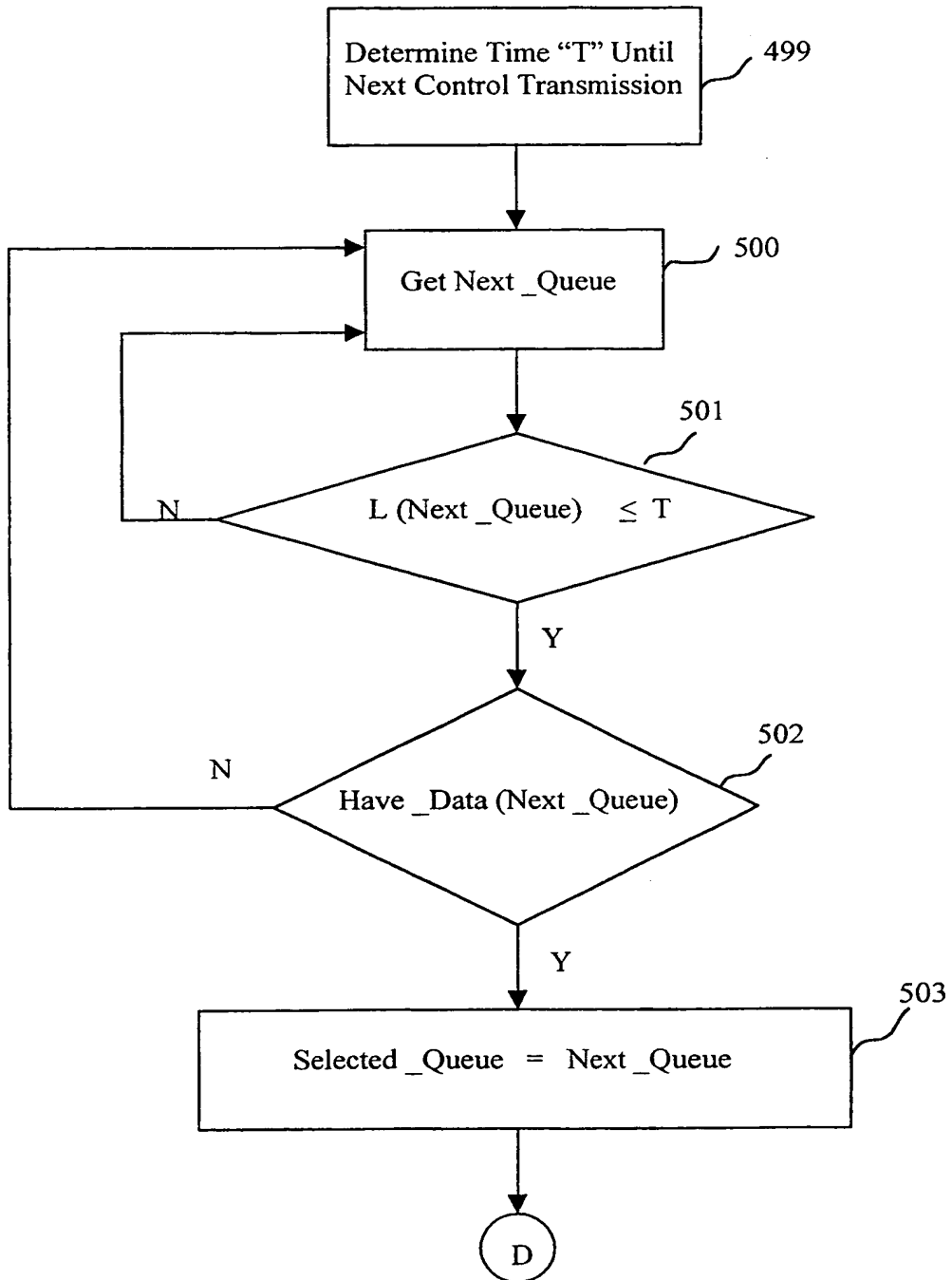


FIGURE 7a

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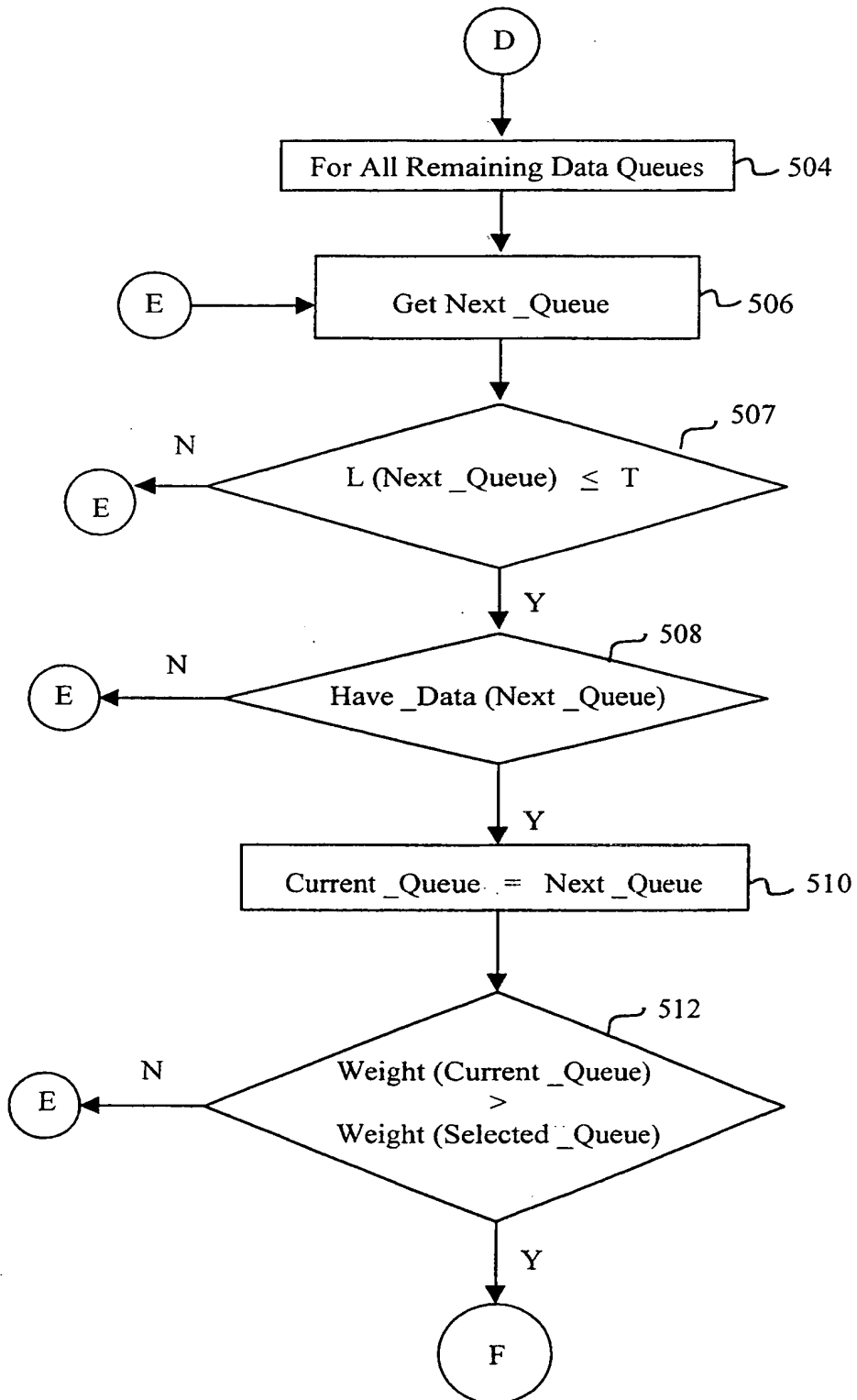


FIGURE 7b

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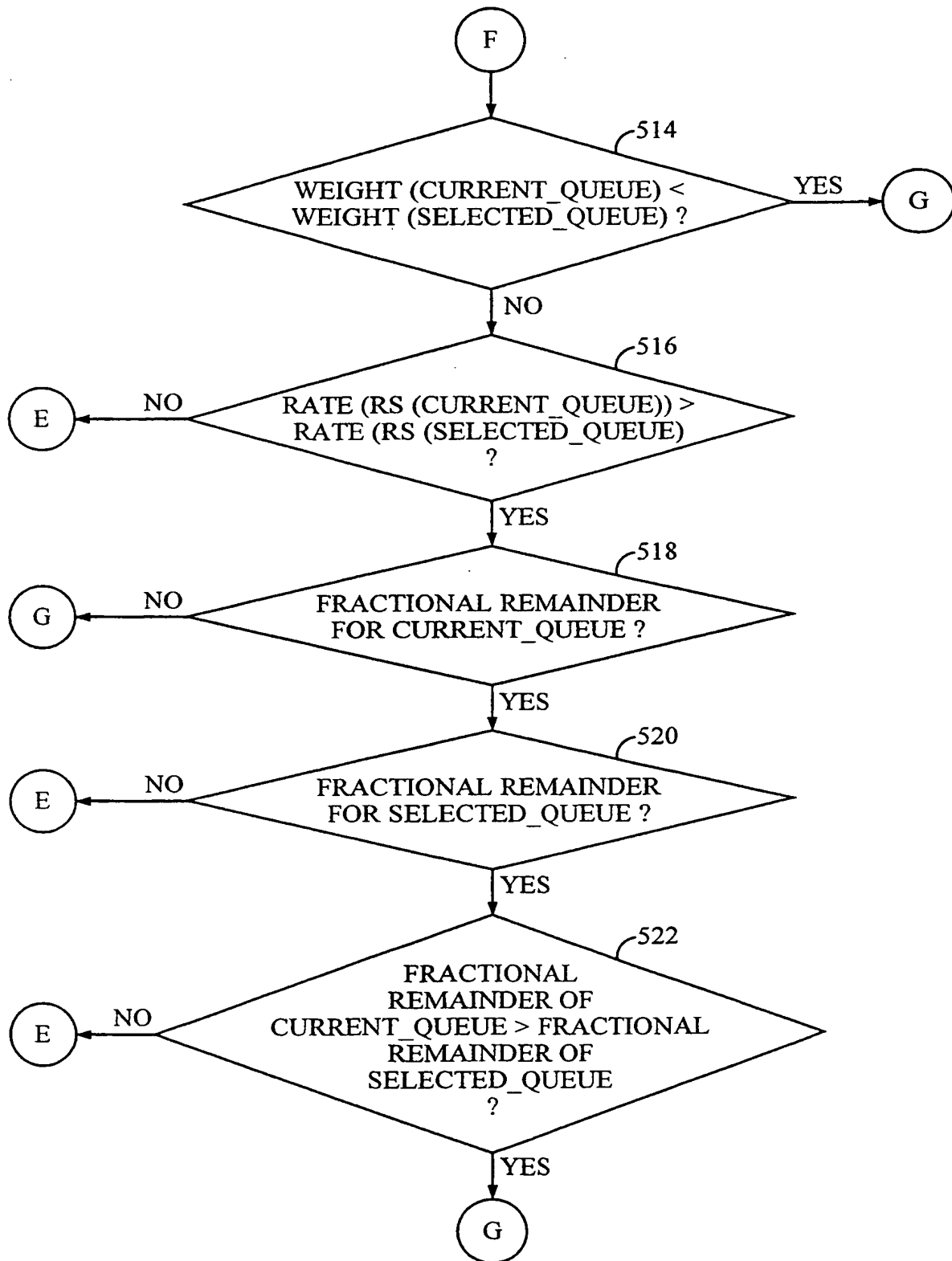


FIG. 7C

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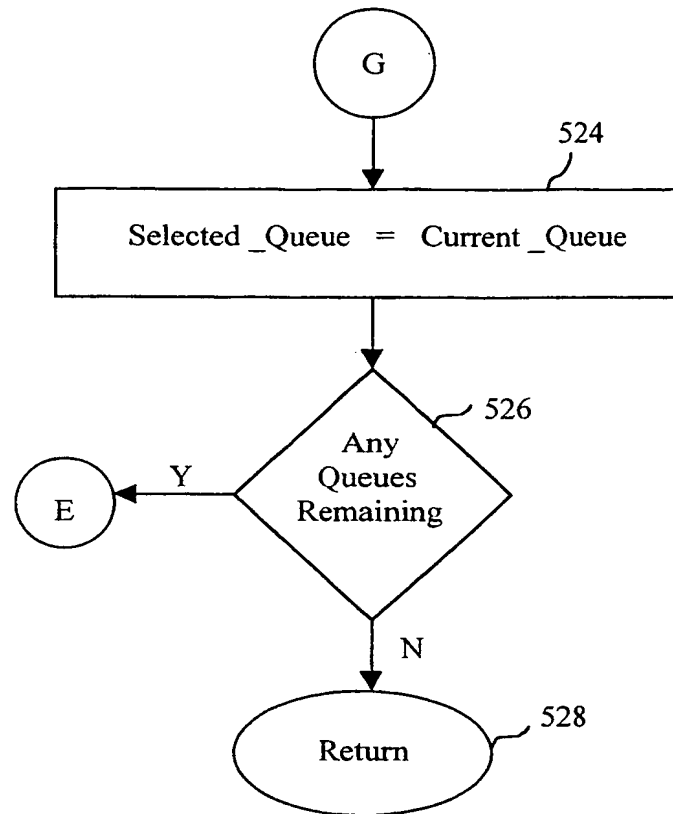


FIGURE 7d

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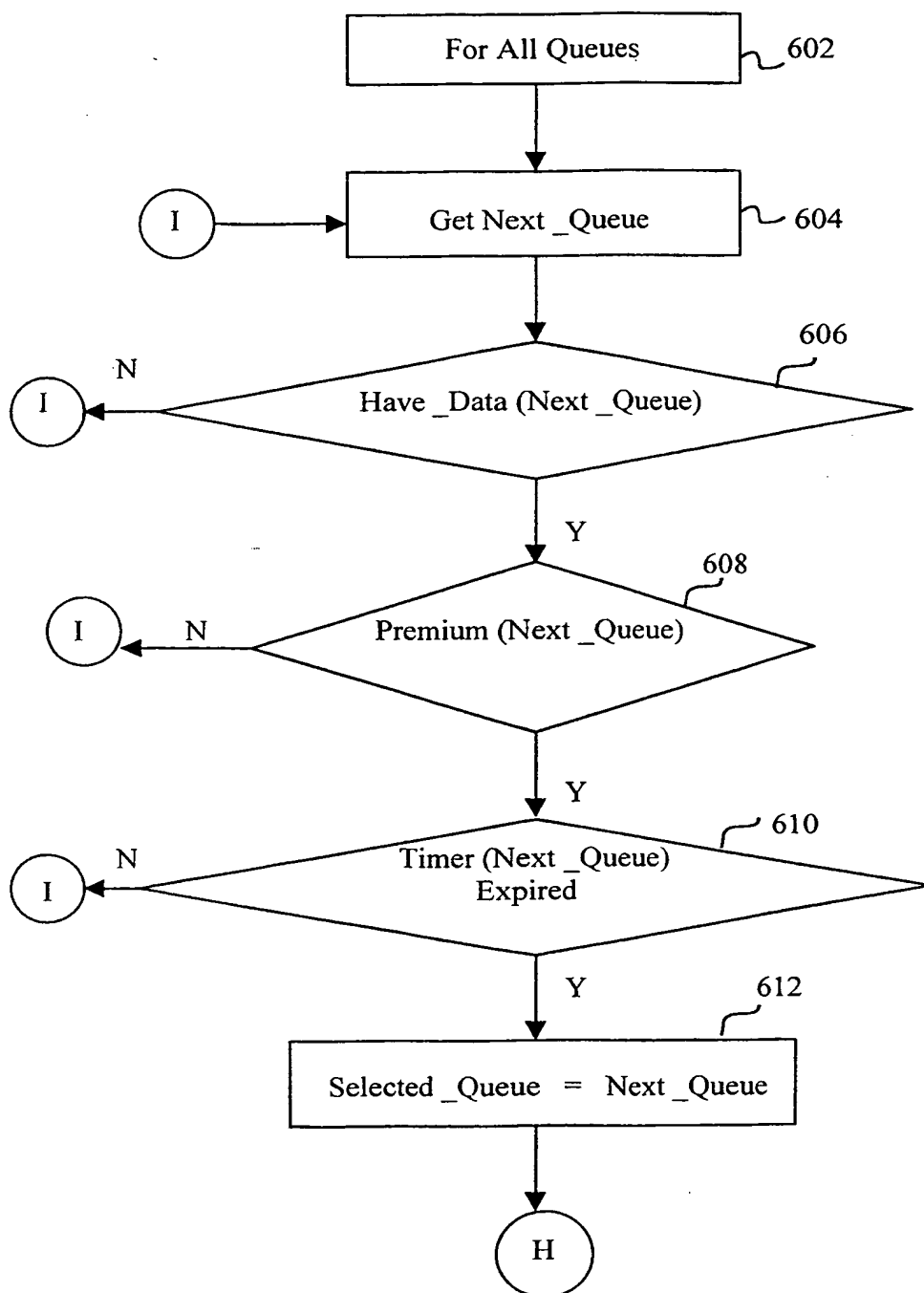


FIGURE 8a

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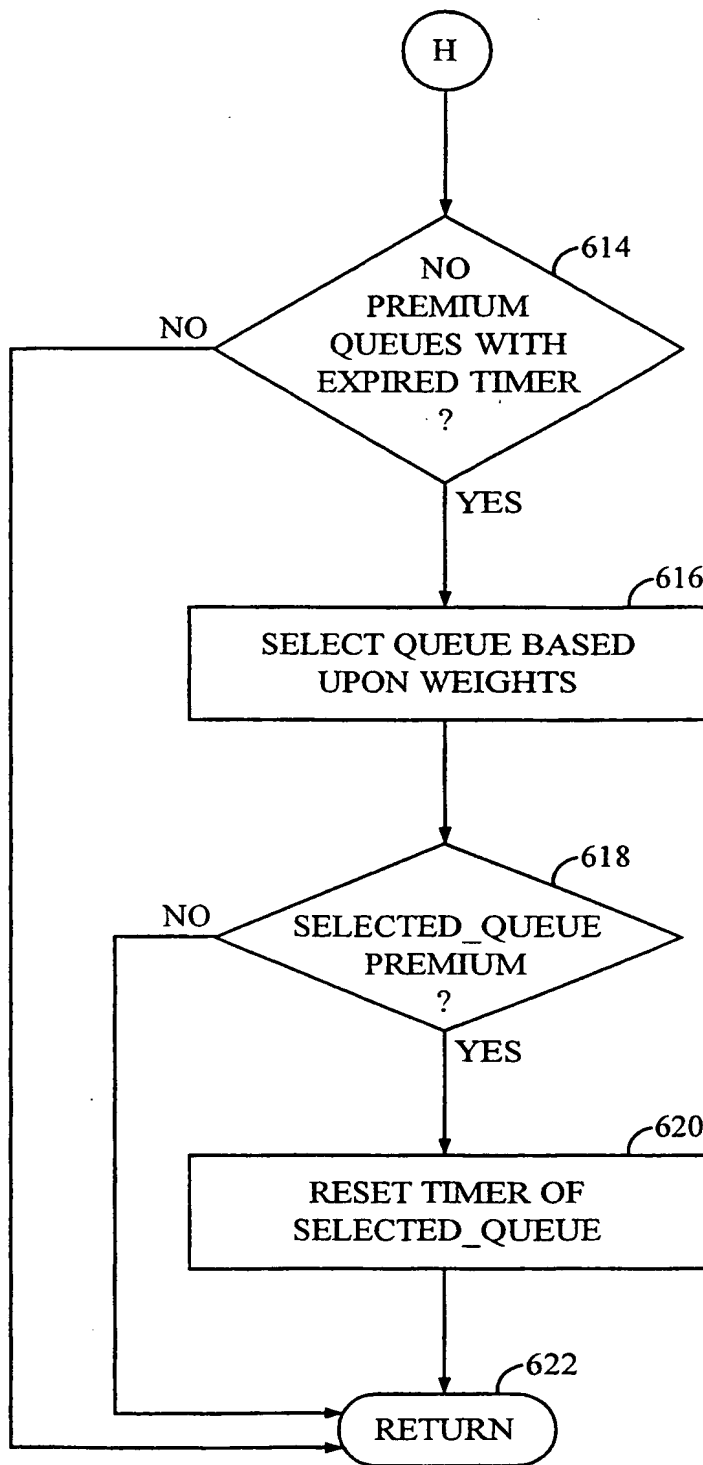


FIG. 8B



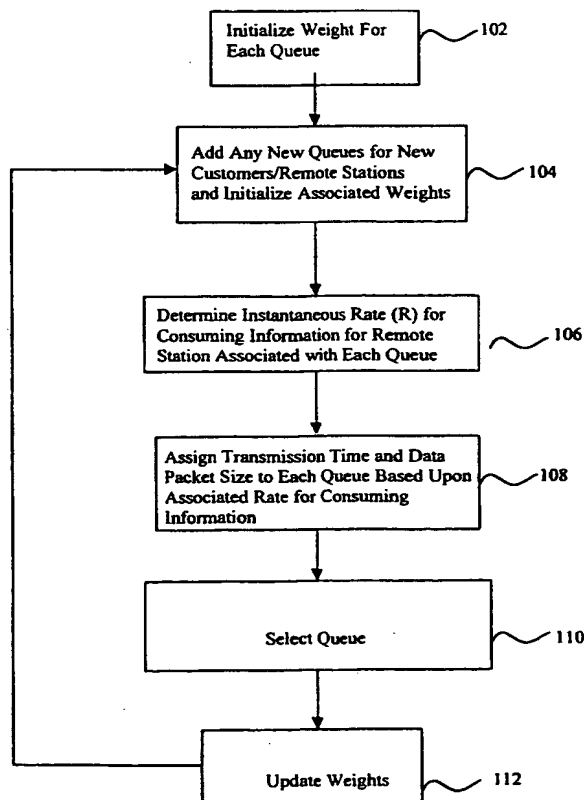
INTERNATIONAL APPLICATION PUBLISHED UNDER THE PATENT COOPERATION TREATY (PCT)

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			(43) International Publication Date: 20 July 2000 (20.07.00)
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(71) Applicant: QUALCOMM INCORPORATED [US/US]; 5775 Morehouse Drive, San Diego, CA 92121-1714 (US).			
(72) Inventors: PANJAK, Rajesh; 7356 Park Village Road, San Diego, CA 92129 (US). SINDHUSHAYANA, Nagabhushana, T.; 10615 Dabney Drive #19, San Diego, CA 92126 (US).		Published With international search report.	
(74) Agents: OGROD, Gregory, D. et al.; Qualcomm Incorporated, 5755 Morehouse Drive, San Diego, CA 92121-1714 (US).		(88) Date of publication of the international search report: 2 November 2000 (02.11.00)	

(54) Title: SYSTEM FOR ALLOCATING RESOURCES IN A COMMUNICATION SYSTEM

(57) Abstract

A communication network having a plurality of subscriber units (6) receive a finite resource from a common node is disclosed. Individual subscriber units (6) may seize the finite resource of the common node to the exclusion of all other subscriber units (6) in the network. A scheduler (12) allocates the finite resource to the individual subscriber units (6) based upon a weight associated with the individual subscriber units (6). The scheduler (12) determines the weight for each of the subscriber units (6) based upon an instantaneous rate of consuming the finite resource.



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INTERNATIONAL SEARCH REPORT

International Application No

PCT/US 00/00801

A. CLASSIFICATION OF SUBJECT MATTER

IPC 7 H04L12/28 H04L12/56 H04Q7/38

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

IPC 7 H04L H04Q

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

WPI Data, PAJ, EPO-Internal, INSPEC

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
Y	WO 98 35514 A (QUALCOMM INC) 13 August 1998 (1998-08-13) abstract page 1, line 8 -page 8, line 15 page 9, line 30 -page 13, line 6	1,2,4, 11,12,14
A	---	3,5,9, 10,13, 15,19,20
	-/--	



Further documents are listed in the continuation of box C.



Patent family members are listed in annex.

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7 July 2000

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INTERNATIONAL SEARCH REPORT

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PCT/US 00/00801

C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
Y	WO 98 45966 A (QUALCOMM INC) 15 October 1998 (1998-10-15) abstract page 1, line 8 -page 7, line 31 page 13, line 12 -page 14, line 2 page 17, line 37 -page 19, line 30 page 20, line 10 - line 25 page 22, line 3 - line 25 page 45, line 10 -page 47, line 33 page 48, line 12 - line 20	1,2,4, 11,12,14
A	---	3,5,9, 10,13, 15,19,20
Y	DOUGLAS C. SCHMIDT, DAVID L.LEVINE, SUMEDH MUNGEE: "The design of the TAO real-time object request broker" ELSEVIER COMPUTER COMMUNICATIONS, no. 21, 1998, pages 1-31, XP002142124 paragraph '001.! paragraph '003.! - paragraph '3.4.7.!	1,2,4, 11,12,14
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A	WO 96 37081 A (ROKE MANOR RESEARCH ;DAVIS SIMON PAUL (GB)) 21 November 1996 (1996-11-21) page 1, line 1 -page 8, line 3 page 9, line 7 -page 11, line 2 page 13, line 18 -page 17, line 24 page 18, line 21 -page 19, line 14 -----	1,11

INTERNATIONAL SEARCH REPORT

information on patent family members

Int. onal Application No

PCT/US 00/00801

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